

# Towards some RandNLA Techniques for Determinant Approximation, Sparse PCA and Analysis of Krylov Subspace Methods

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## LogDet Problem

Given an SPD matrix  $A \in \mathbb{R}^{n \times n}$ , compute (exactly or approximately)  $\log \det(A)$ .

### Additive Error Approximation

Let  $A \in \mathbb{R}^{n \times n}$  be an SPD matrix. For any  $\alpha$  with  $\lambda_1(A) < \alpha$ , define  $B = A/\alpha$  and  $C = I_n - B$ . Then,

$$\log \det(A) = n \log(\alpha) - \sum_{k=1}^{\infty} \frac{\text{Tr}(\log(C^k))}{k}.$$

#### Algorithm 1

**Input:**  $A \in \mathbb{R}^{n \times n}$ , accuracy parameter  $\epsilon > 0$ , integer  $m > 0$ .

- 1 Compute an estimate to the largest eigenvalue of  $A$ ,  $\lambda_1(\tilde{A})$ , using the Power Method.
- 2  $C = I_n - A/(7\lambda_1(\tilde{A}))$
- 3 Create  $p = \lceil 20 \log(2/\delta)/\epsilon^2 \rceil$  i.i.d random Gaussian vectors,  $\mathbf{g}_1, \mathbf{g}_2, \dots, \mathbf{g}_p$ .
- 4 Estimate  $\sum_{k=1}^{\infty} \frac{\text{Tr}(\log(C^k))}{k}$  with a truncated Taylor Series type randomized trace estimator that computes  $\sum_{k=1}^m \left( \frac{1}{p} \sum_{i=1}^p \mathbf{g}_i^T C^k \mathbf{g}_i \right)$

Let  $\widehat{\log \det(A)}$  be the log det approximation of the above procedure. Then, we **prove** that with probability at least  $1 - 2\delta$ ,

$$|\widehat{\log \det(A)} - \log \det(A)| \leq 2\epsilon\Gamma$$

where  $\Gamma = \sum_{i=1}^n \log\left(7 \cdot \frac{\lambda_i(A)}{\lambda_1(A)}\right)$  and  $m \geq \lceil 7\kappa(A) \log(\frac{1}{\epsilon}) \rceil$ .

### Relative Error Approximation

Let  $A \in \mathbb{R}^{n \times n}$  be an SPD matrix whose eigenvalues lie in the interval  $(\theta_1, 1)$ , for some  $0 < \theta_1 < 1$ . Let  $C = I_n - A$ . Then,

$$\log \det(A) = - \sum_{k=1}^{\infty} \frac{\text{Tr}(\log(C^k))}{k}.$$

#### Algorithm 2

**Input:**  $A \in \mathbb{R}^{n \times n}$ , accuracy parameter  $\epsilon > 0$ , integer  $m > 0$ .

- 1  $C = I_n - A$
- 2 Create  $p = \lceil 20 \log(2/\delta)/\epsilon^2 \rceil$  i.i.d random Gaussian vectors,  $\mathbf{g}_1, \mathbf{g}_2, \dots, \mathbf{g}_p$ .
- 3 Estimate  $\sum_{k=1}^{\infty} \frac{\text{Tr}(\log(C^k))}{k}$  with a truncated Taylor Series type randomized trace estimator that computes  $\sum_{k=1}^m \left( \frac{1}{p} \sum_{i=1}^p \mathbf{g}_i^T C^k \mathbf{g}_i \right)$

Let  $\widehat{\log \det(A)}$  be the log det approximation of the above procedure on inputs  $A$  and  $\epsilon$ . Then, we **prove** that with probability at least  $1 - \delta$ ,

$$|\widehat{\log \det(A)} - \log \det(A)| \leq 2\epsilon \cdot |\log \det(A)|.$$

### Citation

C. Boutsidis, P. Drineas, P. Kambadur, E. Kontopoulou, A. Zouzias (2016), *A Randomized Algorithm for Approximating the Log Determinant of a Symmetric Positive Definite Matrix*, under review at Journal of Linear Algebra and its Applications.

ArXiv: <https://arxiv.org/abs/1503.00374>

## Sparse PCA

Given a centered matrix  $X \in \mathbb{R}^{m \times n}$  (the mean of its columns is zero), we seek for a vector  $\mathbf{w}_{opt}$  that solves the optimization problem:

$$\begin{aligned} & \underset{\mathbf{w}}{\text{maximize}} \quad \mathbf{w}^T X^T X \mathbf{w} \\ & \text{subject to} \quad \|\mathbf{w}\|_0 \leq k, \|\mathbf{w}\|_2 \leq 1, \mathbf{w} \in \mathbb{R}^n. \end{aligned}$$

This problem is **NP-hard**  $\rightarrow$  relax to a **problem with convex constraints (but non-convex objective)**:

$$\begin{aligned} & \underset{\mathbf{w}}{\text{maximize}} \quad \mathbf{w}^T X^T X \mathbf{w} \\ & \text{subject to} \quad \|\mathbf{w}\|_1 \leq \sqrt{k}, \|\mathbf{w}\|_2 \leq 1, \mathbf{w} \in \mathbb{R}^n. \end{aligned}$$

#### Algorithm

**Phase 1:** Compute a stationary point  $\tilde{\mathbf{w}}_{opt}$

- 1 Compute the gradient and make a gradient step.
- 2 Project onto the  $l_1$  ball with radius  $\sqrt{k}$  ( $\|\mathbf{w}\|_1$ ).
- 3 Repeat until a threshold for the relative error is exceeded.

**Phase 2:** Invoke a **randomized rounding strategy**.

- 1 Create a Bernoulli distribution and randomly round the entries of  $\mathbf{w}$ .
- 2 Repeat the experiment 10 times and keep the best sparsification.

We prove the following:

Let  $\mathbf{w}_{opt}$  be the optimal solution of the Sparse PCA problem (1) satisfying  $\|\mathbf{w}_{opt}\|_2 = 1$  and  $\|\mathbf{w}_{opt}\|_0 \leq k$ . Let  $\hat{\mathbf{w}}_{opt}$  be the vector returned when the rounding sparsification strategy is applied on the optimal solution  $\tilde{\mathbf{w}}_{opt}$  of the optimization problem (1), with  $s = 200k/\epsilon^2$ , where  $\epsilon \in (0, 1]$  is an accuracy parameter. Then,  $\hat{\mathbf{w}}_{opt}$  has the following properties:

- 1  $\mathbb{E}\|\hat{\mathbf{w}}_{opt}\|_0 \leq s$ .
- 2 With probability at least 3/4,  $\|\hat{\mathbf{w}}_{opt}\|_2 \leq 1 + 0.15\epsilon$ .
- 3 With probability at least 3/4,

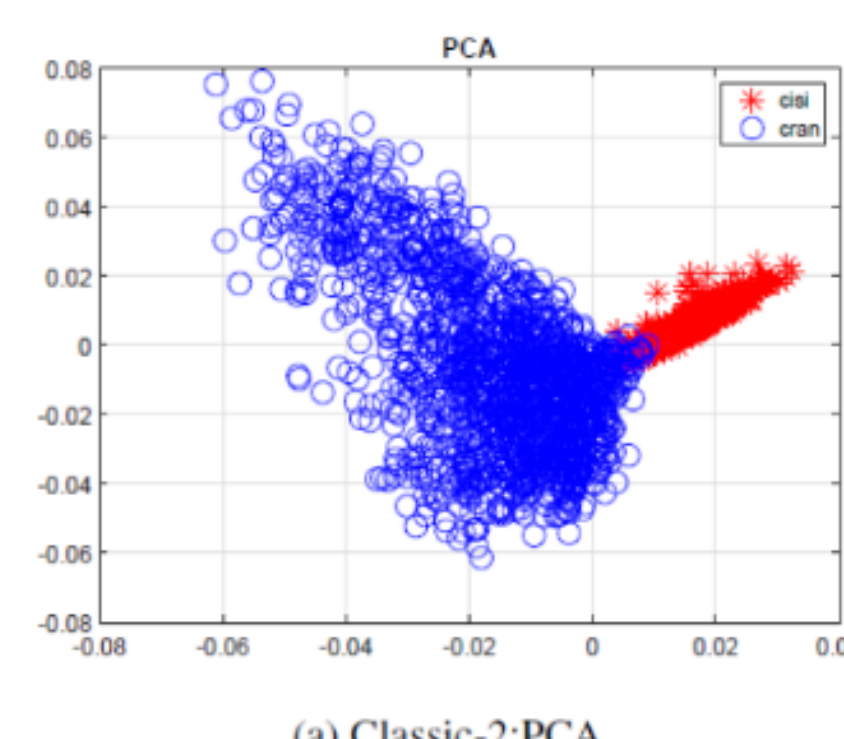
$$\hat{\mathbf{w}}_{opt}^T A \hat{\mathbf{w}}_{opt} \geq \mathbf{w}_{opt}^T A \mathbf{w}_{opt} - \epsilon.$$

### Citation

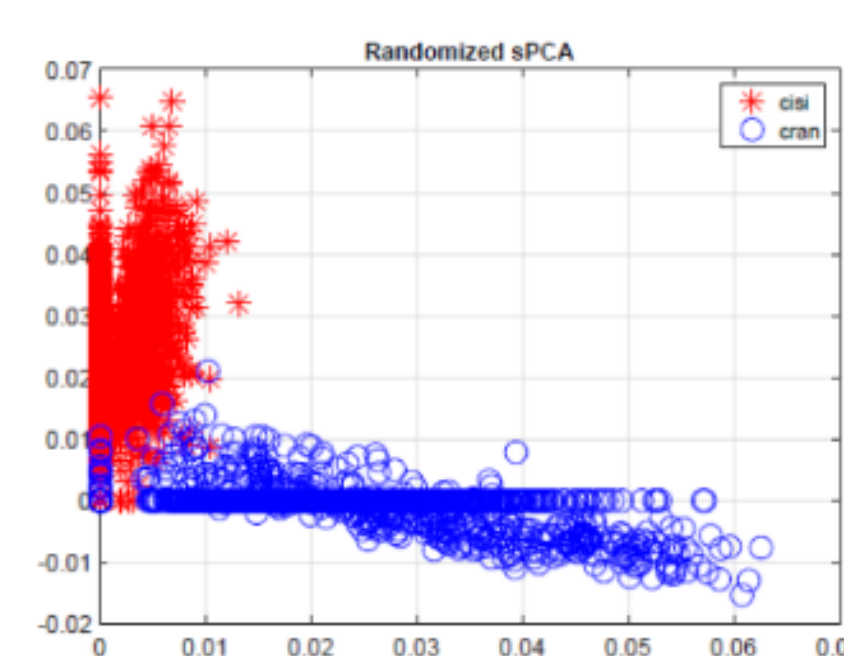
K. Fountoulakis, A. Kundu, E. Kontopoulou, P. Drineas (2016), *A Randomized Rounding Algorithm for Sparse PCA*, under review at ACM Transactions on Knowledge Discovery from Data.

ArXiv: <https://arxiv.org/abs/1508.03337>

## Experiments



(a) Classic-2: PCA.



(b) Classic-2: randomized sPCA.

Datasets  

- Synthetic:  $m = 2^7, n = 2^{12}$
- Classic-2:  $m = 2,858$  documents,  $n = 12,427$  terms
  - CISI collection (1,460 information retrieval abstracts)
  - CRANFIELD collection (1,398 aeronautical systems abstracts)

Table 1: Variance captured by the principal components

	pca	cvx	rsPCA
Top Principal Comp.	41.9%	41.8%	40.8%
Top two Principal Comp.	61.7%	61.5%	59.7%

Table 2: Sparsity of the principal components

	cvx	rsPCA
Top Principal Comp.	65%	91%
2nd Top Principal Comp.	9%	89%

## Krylov Methods

Given a matrix  $A \in \mathbb{R}^{m \times n}$  and a starting guess matrix  $X \in \mathbb{R}^{n \times s}$ , we want to use the block Krylov space  $\mathcal{K}_q(AA^T, AX)$  to approximate the **left singular vector space** of  $A$ .

**We prove:**

- Spectral & Frobenius bounds for the distance between the approximate and the actual space.
- Quality measurements of the bounds relative to the best low-rank approximation.

### Theorem

Let  $\phi(x)$  be a polynomial of degree  $2q + 1$  with odd powers only, such that  $\phi(\Sigma_k)$  is nonsingular. If  $\text{rank}(V_k^T X) = k$  then

$$\|\sin \Theta(\mathcal{K}_q, U_k)\|_{2,F} \leq \|\phi(\Sigma_{k,\perp})\|_2 \|\phi(\Sigma_k)^{-1}\|_2 \|V_{k,\perp}^T X (V_k^T X)^\dagger\|_{2,F}.$$

If, in addition,  $X$  has orthonormal or linearly independent columns, then

$$\|V_{k,\perp}^T X (V_k^T X)^\dagger\|_{2,F} = \|\tan \Theta(X, V_k)\|_{2,F}$$

and

$$\|\sin \Theta(\mathcal{K}_q, U_k)\|_{2,F} \leq \|\phi(\Sigma_{k,\perp})\|_2 \|\phi(\Sigma_k)^{-1}\|_2 \|\tan \Theta(X, V_k)\|_{2,F}.$$

where  $\Theta(\mathcal{K}_q, U_k) \in \mathbb{R}^{k \times k}$  is the diagonal matrix of principal angles between  $\mathcal{K}_q$  and  $\text{range}(U_k)$ .

### Theorem

Let  $\phi(x)$  be a polynomial of degree  $2q + 1$  with odd powers only, such that  $\phi(\Sigma_k)$  is nonsingular and  $\phi(\sigma_i) \geq \sigma_i$ , for  $1 \leq i \leq k$ . If  $\text{rank}(V_k^T X) = k$  then for  $1 \leq i \leq k$ ,

$$\|A - \hat{U}_i \hat{U}_i^T A\|_F \leq \|A - A_i\|_F + \Delta$$

$$\|A - \hat{U}_i \hat{U}_i^T A\|_2 \leq \|A - A_i\|_2 + \Delta$$

$$\sigma_i - \Delta \leq \|\hat{u}_i^T A\|_2 \leq \sigma_i.$$

If, in addition,  $X$  has orthonormal columns, then:

$$\Delta = \|\phi(\Sigma_{k,\perp})\|_2 \|\tan \Theta(X, V_k)\|_F$$

### Citation

I. Ipsen, P. Drineas, E. Kontopoulou, M. Magdon-Ismail (2016), *Structural Convergence Results for Low-Rank Approximations from Block Krylov Spaces*, submitted to SIAM Journal on Matrix Analysis and Applications.

## LogDet Experiments

name	$n$	$nnz$	area of origin
thermal2	1228045	8580313	Thermal
ecology2	999999	4995991	2D/3D
ldoor	952203	42493817	Structural
thermomech_TC	102158	711558	Thermal
boneS01	127224	5516602	Model reduction

exact	logdet(A)		time (sec)		$m$
	approx mean	std	exact	approx mean	
1.3869e6	1.3928e6	964.79	31.28	31.24	149
3.3943e6	3.403e6	1212.8	18.5	10.47	125
1.4429e7	1.4445e7	1683.5	117.91	17.60	33
-546787	-546829.4	553.12	57.84	2.58	77
1.1093e6	1.106e6	247.14	130.4	8.48	125